PATENT

Filed: January 26, 2001

CASE NO.: ARC920010006US1

Serial No.: 09/771,239

December 7, 2004

Page 2

A method for identifying or disabling at least one traitor receiver with (currently amended) 1.

at least one associated unique, compromised decryption key in a broadcast encryption system, comprising:

receiving a set of subsets derived from a tree defining leaves, each leaf representing a

respective receiver;

identifying at least one traitor subset from the set of subsets as containing at least one leaf

representing a candidate traitor receiver;

using the traitor subset, identifying or disabling the traitor receiver; and

determining whether the traitor subset represents at least two traitor receivers, and if so,

dividing the traitor subset into two child sets.

2. (canceled).

3. (previously presented) The method of Claim 1, further comprising determining whether the

traitor subset is a member of a frontier set, and if so, removing a complementary subset from the frontier

set.

(original) The method of Claim 1, wherein the act of identifying or disabling includes 4.

encoding plural subsets of the set of subsets with a false key.

5. (original) The method of Claim 4, further comprising executing a binary search on the set

of subsets using probabilities.

(TUE) DEC 7 2004 15:51/ST. 15:50/No. 6833031378 P 3

FROM ROGITZ 619 338 8078

3 %

CASE NO.: ARC920010006US1

Serial No.: 09/771,239

December 7, 2004

Page 3

PATENT Filed: January 26, 2001

6. (original) The method of Claim 5, wherein the binary search ends by determining that the

difference between a probability p_j of decrypting a message when the first j subsets contain the false key and

a probability $p_{j,i}$ of decrypting a message when the first j-1 subsets contain the false key is at least equal to

a predetermined probability.

7. (original) The method of Claim 6, wherein the traitor subset is identified when $|p_{j_1}-p_{j_1}|$

> p/m, wherein m is the number of subsets in the set of subsets.

8. (original) The method of Claim 1, wherein the set of subsets is generated by:

assigning each receiver in a group of receivers respective private information I,;

selecting at least one session encryption key K;

partitioning receivers not in a revoked set R into a set of disjoint subsets S₁₁,...,S_{im} having

associated subset keys L₁₁,...,L_{lm}; and

encrypting the session key K and the false key with the subset keys $L_{i_1},...,L_{i_m}$.

9. (original) The method of Claim 8, wherein the tree includes a root and plural nodes, each

node having an associated key, and wherein each receiver is assigned keys from all nodes in a direct path

between a leaf representing the receiver and the root.

CASE NO.: ARC920010006US1

Serial No.: 09/771,239

December 7, 2004

Page 4

PATENT Filed: January 26, 2001

10. (original) The method of Claim 8, wherein the tree includes a root and plural nodes, each node associated with a set of labels, and wherein each receiver is assigned labels from all nodes hanging from

a direct path between the receiver and the root but not from nodes in the direct path.

11. (original) The method of Claim 10, wherein the revoked set R defines a spanning tree, and

wherein the method includes:

initializing a cover tree T as the spanning tree;

iteratively removing nodes from the cover tree T and adding nodes to the cover tree T until

the cover tree T has at most one node.

12. (original) A computer program device, comprising:

a computer program storage device including a program of instructions usable by a computer,

comprising:

logic means for accessing a tree to generate a set of subsets of the tree, the tree including

leaves representing at least one traitor device characterized by a compromised key;

logic means for encrypting a false key j times and for encrypting a session key m-j times,

wherein m is a number of subsets in the set of subsets;

logic means responsive to the means for encrypting for identifying a traitor subset; and

logic means for using the traitor subset to identify or disable the traitor device.

13. (currently amended) The computer program device of Claim 12, further comprising:

(TUE) DEC 7 2004 15:51/ST. 15:50/No. 6833031378 P 5

FROM ROGITZ 619 338 8078

CASE NO.: ARC920010006US1

Serial No.: 09/771,239

December 7, 2004

Page 5

PATENT Filed: January 26, 2001

logic means for determining whether the traitor subset represents at least two candidate traitor

devices, and if so, dividing the traitor subset into two child sets.

14. (original) The computer program device of Claim 13, further comprising logic means for

determining whether the traitor subset is a member of a frontier set, and if so, removing a complementary

subset from the frontier set.

15. (original) The computer program device of Claim 12, further comprising logic means for

executing a binary search on the set of subsets using probabilities.

16. (original) The computer program device of Claim 15, wherein the binary search ends by

determining that the difference between a probability p_i of decrypting a message when the first j subsets

contain the false key and a probability p_{j-1} of decrypting a message when the first j-1 subsets contain the false

key is at least equal to a predetermined probability.

17. (original) The computer program device of Claim 16, wherein the traitor subset is identified

when $|p_{i,1}-p_{i,1}| > p/m$, wherein m is the number of subsets in the set of subsets.

18. (original) The method of Claim 12, wherein the set of subsets is generated by logic means

including:

U

CASE NO.: ARC920010006US1

Serial No.: 09/771,239

December 7, 2004

Page 6

PATENT Filed: January 26, 2001

logic means for assigning each receiver in a group of receivers respective private information

Iu;

logic means for selecting at least one session encryption key K;

logic means for partitioning receivers not in a revoked set R into a set of disjoint subsets

 $S_{i1},...S_{im}$ having associated subset keys $L_{i1},...,L_{im}$; and

logic means for encrypting the session key K and the false key with the subset keys

 $L_{i1},...,L_{im}$

19. (original) The computer program device of Claim 18, wherein the tree includes a root and

plural nodes, each node having an associated key, and wherein each receiver is assigned keys from all nodes

hanging from a direct path between the receiver and the root but not from nodes in the direct path.

20. (original) A computer programmed with instructions to cause the computer to execute method

acts including:

using a false key to encode plural subsets representing stateless receivers, at least one traitor

receiver of which is associated with at least one compromised key that has been obtained by at least

one pirate receiver; and

using the pirate receiver or a clone thereof, determining the identity of the traitor receiver,

or rendering the pirate receiver or clone thereof useless for decrypting data using the compromised

key.

CASE NO.: ARC920010006U\$1

Serial No.: 09/771,239 December 7, 2004

Page 7

PATENT Filed: January 26, 2001

21. (original) The computer of Claim 20, wherein the subsets define a set of subsets, and the

method acts undertaken by the computer further include:

receiving the set of subsets derived from a tree defining leaves, each leaf representing a

respective receiver;

identifying at least one traitor subset from the set of subsets as containing at least one leaf

representing the traitor receiver; and

using the traitor subset, identifying the traitor receiver.

22. (currently amended) The computer of Claim 21, wherein the method acts undertaken by the

computer further comprise:

determining whether the traitor subset represents at least two candidate traitor receivers, and

if so, dividing the traitor subset into two child sets.

(original) The computer of Claim 22, wherein the method acts undertaken by the computer 23.

further comprise determining whether the traitor subset is a member of a frontier set, and if so, removing

a complementary subset from the frontier set.

(original) The computer of Claim 21, wherein the act of identifying includes: 24.

encoding plural subsets of the set of subsets with the false key.

CASE NO.: ARC920010006US1

Serial No.: 09/771,239 December 7, 2004

Page 8

PATENT Filed: January 26, 2001

25. (original) The computer of Claim 24, wherein the method acts undertaken by the computer

further comprise executing a binary search on the set of subsets using probabilities.

26. (original) The computer of Claim 25, wherein the binary search ends by determining that a

probability p_j of decrypting a message when the first j subsets contain the false key is at least equal to a

predetermined probability.

(original) The computer of Claim 26, wherein the traitor subset is identified when $|p_{i,1}-p_{i,j}|$ 27.

> p/m, wherein m is the number of subsets in the set of subsets.

28. (original) The computer of Claim 21, wherein the set of subsets is generated by:

assigning each receiver in a group of receivers respective private information I,;

selecting at least one session encryption key K;

partitioning receivers not in a revoked set R into a set of disjoint subsets S₁₁,...S_{lm} having

associated subset keys Lii,...,Lin; and

encrypting the session key K and the false key with the subset keys $L_1, ..., L_m$, wherein the

tree includes a root and plural nodes, each node being associated with a set of labels, and wherein

each receiver is assigned labels from all nodes hanging from a direct path between the receiver and

the root but not from nodes in the direct path.

CASE NO.: ARC920010006US1

Serial No.: 09/771,239 December 7, 2004

Page 9

Filed: January 26, 2001

- 29. (original) The method of Claim 1, further comprising identifying or disabling plural traitor receivers embodied in a clone.
- 30. (original) The method of Claim 1, wherein the act of identifying or disabling includes encoding the first j subsets of the set of subsets with a false key.

1053-122-AM1